

# A Demonstration of a Relaying Selection Scheme for Maximizing a Diamond Network's Throughput

## INTRODUCTION

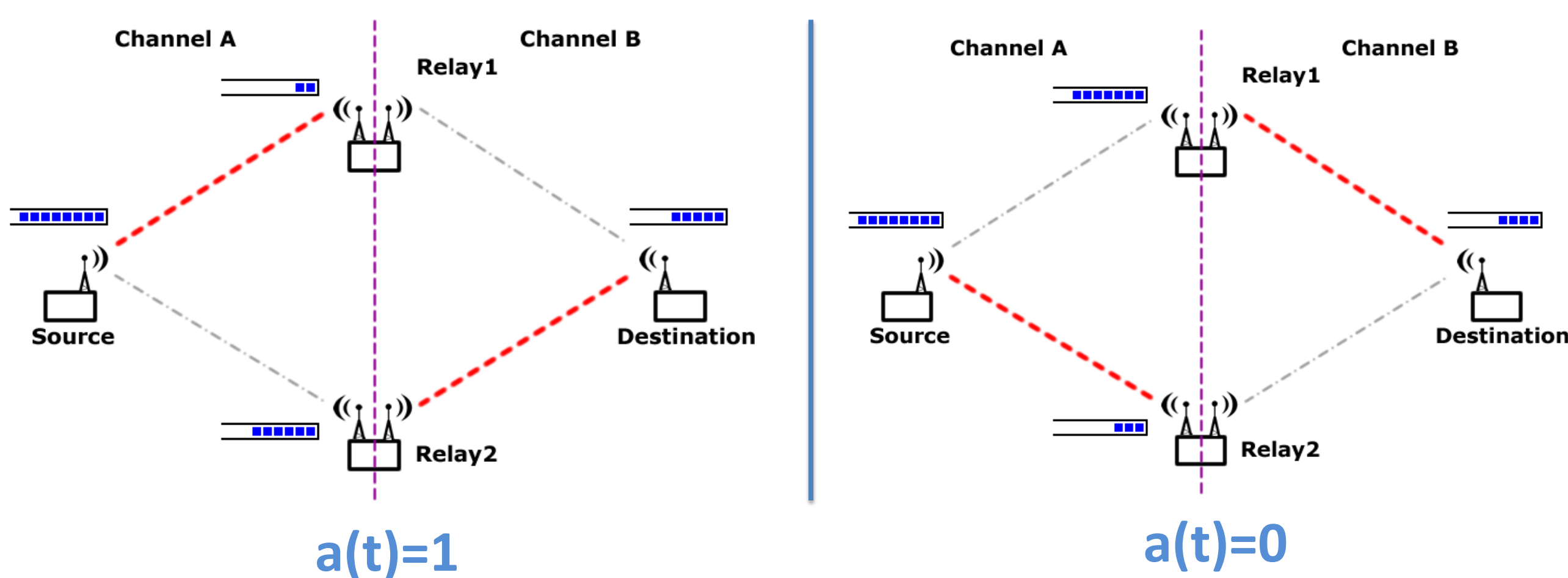
**Aim:** Maximize throughput while achieving network stability.

**Goal:** Select schedules in the Diamond Relay Network, towards maximizing the total network's throughput.

**Method:** Rely on optimization theory tools and Lyapunov drift to obtain optimal schedules. A network controller  $a(t)$ , chooses the optimal between two feasible scheduling activation sets.

**Implementation:** Per packet-level configuration using Click Modular Router and Ath9k driver.

**Access Method:** A TDMA framing over WiFi .



Red lines denote the activation of two feasible scheduling sets by controller  $a(t)$ .

## IMPLEMENTATION ISSUES

- ✓ Operating in a single frequency with CSMA prevents us from enabling parallel transmissions i.e.  $S \rightarrow R1$  and  $R2 \rightarrow D$  without collisions.
- ✓ *Solution:* Use of two different channels operating in each hop, in order to enable independent schedules.
- ✓ Scheduling decisions are taken in the IP layer rather than the MAC layer, since gathering and handling control data is more flexible with Click Modular Router.

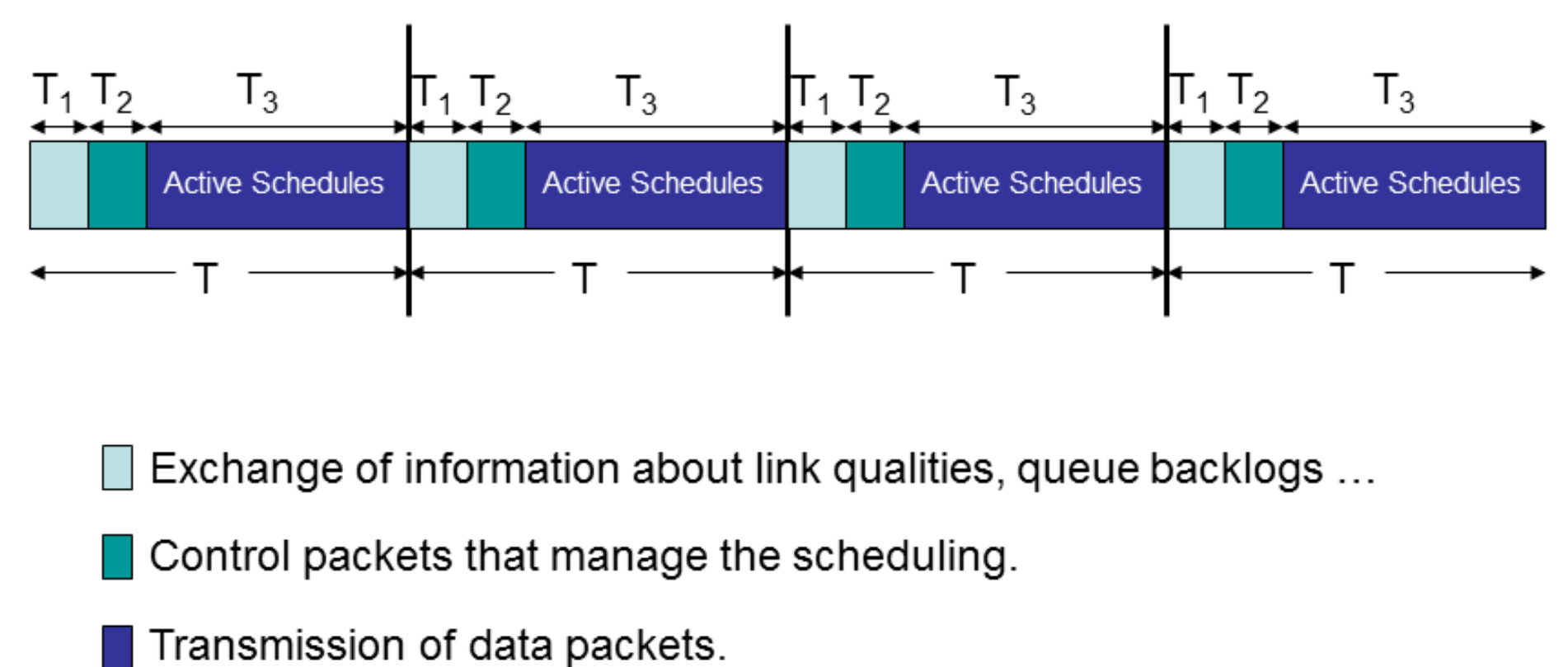
## SYSTEM SETUP

- ✓ 1 Source, 2 Relays, 1 Destination.
- ✓ Each node maintains a backlog data queue  $Q(t)$ .
- ✓ Channel States:  $S(t)$ , Service Rates:  $\mu(t)$ .
- ✓ Lyapunov function:  $L(Q(t)) = \sum_i Q_i(t)^2$
- ✓ Drift Expression:  $\Delta(t) = E[L(Q(t+1)) - L(Q(t)) | Q(t)]$
- ✓ Slotted time, TDMA frame structure.
- ✓ T1 intervals for gathering Network State Information, T2 for reporting schedules and T3 for actual transmission .

## SOLUTION APPROACH

- ✓ Minimize a bound on the drift expression with respect to  $a(t)$

$$\min_{a(t)} \Delta(a(t), t)$$



TDMA frame structure.

## MAX WEIGHT RULE ALGORITHM

- ✓ T1 interval: Source Node gathers Network State Information from its neighbors.
- ✓ T2 interval: Source takes a scheduling decision.
  - If  $\Delta Q_{SR_1}(t)\mu_{SR_1}(t) + Q_{R_2}(t)\mu_{R_2,D}(t) < \Delta Q_{SR_2}(t)\mu_{SR_2}(t) + Q_{R_1}(t)\mu_{R_1,D}(t)$
  - Set  $a(t)=1$ , and transmit over the first feasible set.
  - Otherwise, set  $a(t)=0$ , and transmit over the second feasible set.
- ✓ T3 interval: Transmit over the selected schedule set with rate  $\mu(t)$ .